

# Protothreads

## The Protothreads Library 1.1 Reference Manual

April 2005



Adam Dunkels  
adam@sics.se

Swedish Institute of Computer Science

## Contents

<b>1</b>	<b>The Protothreads Library</b>	<b>1</b>
<b>2</b>	<b>The Protothreads Library 1.1 Module Index</b>	<b>3</b>
<b>3</b>	<b>The Protothreads Library 1.1 File Index</b>	<b>4</b>
<b>4</b>	<b>The Protothreads Library 1.1 Module Documentation</b>	<b>4</b>
<b>5</b>	<b>The Protothreads Library 1.1 File Documentation</b>	<b>23</b>

## 1 The Protothreads Library

### Author:

Adam Dunkels <[adam@sics.se](mailto:adam@sics.se)>

Protothreads are a type of lightweight stackless threads designed for severely memory constrained systems such as deeply embedded systems or sensor network nodes. Protothreads provides linear code execution for event-driven systems implemented in C. Protothreads can be used with or without an RTOS.

Protothreads are an extremely lightweight, stackless type of threads that provides a blocking context on top of an event-driven system, without the overhead of per-thread stacks. The purpose of protothreads is to implement sequential flow of control without complex state machines or full multi-threading. Protothreads provides conditional blocking inside C functions.

Main features:

- No machine specific code - the protothreads library is pure C
- Does not use error-prone functions such as `longjmp()`
- Very small RAM overhead - only two bytes per protothread
- Can be used with or without an OS
- Provides blocking wait without full multi-threading or stack-switching

Examples applications:

- Memory constrained systems
- Event-driven protocol stacks
- Deeply embedded systems
- Sensor network nodes

See also:

[Example programs](#)

[Protothreads API documentation](#)

The protothreads library is released under a BSD-style license that allows for both non-commercial and commercial usage. The only requirement is that credit is given.

More information and new version of the code can be found at the Protothreads homepage:

<http://www.sics.se/~adam/pt/>

## 1.1 Authors

The protothreads library was written by Adam Dunkels <[adam@sics.se](mailto:adam@sics.se)> with support from Oliver Schmidt <[ol.sc@web.de](mailto:ol.sc@web.de)>.

## 1.2 Using protothreads

Using protothreads in a project is easy: simply copy the files [pt.h](#), [lc.h](#) and [lc-switch.h](#) into the include files directory of the project, and `#include "pt.h"` in all files that should use protothreads.

## 1.3 Protothreads

Protothreads are an extremely lightweight, stackless threads that provides a blocking context on top of an event-driven system, without the overhead of per-thread stacks. The purpose of protothreads is to implement sequential flow of control without using complex state machines or full multi-threading. Protothreads provides conditional blocking inside a C function.

In memory constrained systems, such as deeply embedded systems, traditional multi-threading may have a too large memory overhead. In traditional multi-threading, each thread requires its own stack, that typically is over-provisioned. The stacks may use large parts of the available memory.

The main advantage of protothreads over ordinary threads is that protothreads are very lightweight: a protothread does not require its own stack. Rather, all protothreads run on the same stack and context switching is done by stack rewinding. This is advantageous in memory constrained systems, where a stack for a thread might use a large part of the available memory. A protothread only requires only two bytes of memory per protothread. Moreover, protothreads are implemented in pure C and do not require any machine-specific assembler code.

A protothread runs within a single C function and cannot span over other functions. A protothread may call normal C functions, but cannot block inside a called function. Blocking inside nested function calls is instead made by spawning a separate protothread for each potentially blocking function. The advantage of this approach is that blocking is explicit: the programmer knows exactly which functions that block that which functions the never blocks.

Protothreads are similar to asymmetric co-routines. The main difference is that co-routines uses a separate stack for each co-routine, whereas protothreads are stackless. The most similar mechanism to protothreads are Python generators. These are also stackless constructs, but have a different purpose. Protothreads provides blocking contexts inside a C function, whereas Python generators provide multiple exit points from a generator function.

## 1.4 Local variables

### Note:

Because protothreads do not save the stack context across a blocking call, local variables are not preserved when the protothread blocks. This means that local variables should be used with utmost care - if in doubt, do not use local variables inside a protothread!

## 1.5 Scheduling

A protothread is driven by repeated calls to the function in which the protothread is running. Each time the function is called, the protothread will run until it blocks or exits. Thus the scheduling of protothreads is done by the application that uses protothreads.

## 1.6 Implementation

Protothreads are implemented using local continuations. A local continuation represents the current state of execution at a particular place in the program, but does not provide any call history or local variables. A local continuation can be set in a specific function to capture the state of the function. After a local continuation has been set can be resumed in order to restore the state of the function at the point where the local continuation was set.

Local continuations can be implemented in a variety of ways:

1. by using machine specific assembler code,
2. by using standard C constructs, or
3. by using compiler extensions.

The first way works by saving and restoring the processor state, except for stack pointers, and requires between 16 and 32 bytes of memory per protothread. The exact amount of memory required depends on the architecture.

The standard C implementation requires only two bytes of state per protothread and utilizes the C `switch()` statement in a non-obvious way that is similar to Duff's device. This implementation does, however, impose a slight restriction to the code that uses protothreads: a protothread cannot perform a blocking wait (`PT_WAIT_UNTIL()` or `PT_YIELD()`) inside a `switch()` statement.

Certain compilers has C extensions that can be used to implement protothreads. GCC supports label pointers that can be used for this purpose. With this implementation, protothreads require 4 bytes of RAM per protothread.

## 2 The Protothreads Library 1.1 Module Index

### 2.1 The Protothreads Library 1.1 Modules

Here is a list of all modules:

<b>Protothreads</b>	<b>4</b>
<b>Protothread semaphores</b>	<b>19</b>
<b>Local continuations</b>	<b>22</b>

Examples	10
----------	----

## 3 The Protothreads Library 1.1 File Index

### 3.1 The Protothreads Library 1.1 File List

Here is a list of all documented files with brief descriptions:

<a href="#">lc-addrlabels.h</a> (Implementation of local continuations based on the "Labels as values" feature of gcc )	23
<a href="#">lc-switch.h</a> (Implementation of local continuations based on switch() statment )	24
<a href="#">lc.h</a> (Local continuations )	24
<a href="#">pt-sem.h</a> (Couting semaphores implemented on protothreads )	25
<a href="#">pt.h</a> (Protothreads implementation )	26

## 4 The Protothreads Library 1.1 Module Documentation

### 4.1 Protothreads

#### 4.1.1 Detailed Description

Protothreads are implemented in a single header file, [pt.h](#), which includes the local continuations header file, [lc.h](#). This file in turn includes the actual implementation of local continuations, which typically also is contained in a single header file.

#### Files

- file [pt.h](#)  
*Protothreads implementation.*

#### Modules

- group [Protothread semaphores](#)
- group [Local continuations](#)

#### Defines

- #define [PT\\_THREAD](#)(name\_args)  
*Declaration of a protothread.*
- #define [PT\\_INIT](#)(pt)  
*Initialize a protothread.*

- `#define PT_BEGIN(pt)`  
*Declare the start of a protothread inside the C function implementing the protothread.*
- `#define PT_WAIT_UNTIL(pt, condition)`  
*Block and wait until condition is true.*
- `#define PT_WAIT_WHILE(pt, cond)`  
*Block and wait while condition is true.*
- `#define PT_WAIT_THREAD(pt, thread)`  
*Block and wait until a child protothread completes.*
- `#define PT_SPAWN(pt, child, thread)`  
*Spawn a child protothread and wait until it exits.*
- `#define PT_RESTART(pt)`  
*Restart the protothread.*
- `#define PT_EXIT(pt)`  
*Exit the protothread.*
- `#define PT_END(pt)`  
*Declare the end of a protothread.*
- `#define PT_SCHEDULE(f)`  
*Schedule a protothread.*
- `#define PT_YIELDING()`  
*Declare that a protothread can yield.*
- `#define PT_YIELD(pt)`  
*Yield from the current protothread.*

#### 4.1.2 Define Documentation

##### 4.1.2.1 `#define PT_BEGIN(pt)`

Declare the start of a protothread inside the C function implementing the protothread.

This macro is used to declare the starting point of a protothread. It should be placed at the start of the function in which the protothread runs. All C statements above the `PT_BEGIN()` invocation will be executed each time the protothread is scheduled.

##### Parameters:

*pt* A pointer to the protothread control structure.

Example:

```
PT_THREAD(producer(struct pt *p, int event)) {
    PT_BEGIN(p);
    while(1) {
        PT_WAIT_UNTIL(p, event == CONSUMED || event == DROPPED);
        produce();
        PT_WAIT_UNTIL(p, event == PRODUCED);
    }

    PT_END(p);
}
```

#### 4.1.2.2 #define PT\_END(pt)

Declare the end of a protothread.

This macro is used for declaring that a protothread ends. It should always be used together with a matching [PT\\_BEGIN\(\)](#) macro.

**Parameters:**

*pt* A pointer to the protothread control structure.

#### 4.1.2.3 #define PT\_EXIT(pt)

Exit the protothread.

This macro causes the protothread to exit. If the protothread was spawned by another protothread, the parent protothread will become unblocked and can continue to run.

**Parameters:**

*pt* A pointer to the protothread control structure.

#### 4.1.2.4 #define PT\_INIT(pt)

Initialize a protothread.

Initializes a protothread. Initialization must be done prior to starting to execute the protothread.

**Parameters:**

*pt* A pointer to the protothread control structure.

Example:

```
void main(void) {
    struct pt p;
    int event;

    PT_INIT(&p);
    while(PT_SCHEDULE(consumer(&p, event))) {
        event = get_event();
    }
}
```

**See also:**

[PT\\_SPAWN\(\)](#)

#### 4.1.2.5 #define PT\_RESTART(pt)

Restart the protothread.

This macro will block and cause the running protothread to restart its execution at the place of the `PT_BEGIN()` call.

**Parameters:**

*pt* A pointer to the protothread control structure.

#### 4.1.2.6 #define PT\_SCHEDULE(f)

Schedule a protothread.

This function schedules a protothread. The return value of the function is non-zero if the protothread is running or zero if the protothread has exited.

Example

```
void main(void) {
    struct pt p;
    int event;

    PT_INIT(&p);
    while(PT_SCHEDULE(consumer(&p, event))) {
        event = get_event();
    }
}
```

**Parameters:**

*f* The call to the C function implementing the protothread to be scheduled

#### 4.1.2.7 #define PT\_SPAWN(pt, child, thread)

Spawn a child protothread and wait until it exits.

This macro spawns a child protothread and waits until it exits. The macro can only be used within a protothread.

Example:

```
static struct pt parent_pt, child_pt;
int should_spawn_flag;

PT_THREAD(child(struct pt *pt)) {
    PT_BEGIN(pt);

    while(all_items_processed()) {
        process_item();
        PT_WAIT_UNTIL(pt, item_processed());
    }

    PT_END(pt);
}

PT_THREAD(parent(void)) {
    PT_BEGIN(&parent_pt);

    if(should_spawn_flag) {
        PT_SPAWN(&parent_pt, &child_pt, child(&child_pt));
    }
}
```



```

    }

    PT_END(&parent_pt);
}

```

**Parameters:**

- pt* A pointer to the protothread control structure.
- child* A pointer to the child protothread's control structure.
- thread* The child protothread with arguments

**4.1.2.8 #define PT\_THREAD(name\_args)**

Declaration of a protothread.

This macro is used to declare a protothread. All protothreads must be declared with this macro.

Example:

```

PT_THREAD(consumer(struct pt *p, int event)) {
    PT_BEGIN(p);
    while(1) {
        PT_WAIT_UNTIL(p, event == AVAILABLE);
        consume();
        PT_WAIT_UNTIL(p, event == CONSUMED);
        acknowledge_consumed();
    }
    PT_END(p);
}

```

**Parameters:**

- name\_args* The name and arguments of the C function implementing the protothread.

**4.1.2.9 #define PT\_WAIT\_THREAD(pt, thread)**

Block and wait until a child protothread completes.

This macro schedules a child protothread. The current protothread will block until the child protothread completes.

**Note:**

The child protothread must be manually initialized with the [PT\\_INIT\(\)](#) function before this function is used.

**Parameters:**

- pt* A pointer to the protothread control structure.
- thread* The child protothread with arguments

Example:

```

PT_THREAD(child(struct pt *p, int event)) {
    PT_BEGIN(p);

    PT_WAIT_UNTIL(p, event == EVENT1);

    PT_END(p);
}

```

```
PT_THREAD(parent(struct pt *p, struct pt *child_pt, int event)) {
    PT_BEGIN(p);

    PT_INIT(child_pt);

    PT_WAIT_THREAD(p, child(child_pt, event));

    PT_END(p);
}
```

See also:

[PT\\_SPAWN\(\)](#)

#### 4.1.2.10 #define PT\_WAIT\_UNTIL(pt, condition)

Block and wait until condition is true.

This macro blocks the protothread until the specified condition is true.

**Parameters:**

*pt* A pointer to the protothread control structure.

*condition* The condition.

Example:

```
PT_THREAD(seconds(struct pt *p)) {
    PT_BEGIN(p);

    PT_WAIT_UNTIL(p, time >= 2 * SECOND);
    printf("Two seconds have passed\n");

    PT_END(p);
}
```

#### 4.1.2.11 #define PT\_WAIT\_WHILE(pt, cond)

Block and wait while condition is true.

This function blocks and waits while condition is true. See [PT\\_WAIT\\_UNTIL\(\)](#).

**Parameters:**

*pt* A pointer to the protothread control structure.

*cond* The condition.

#### 4.1.2.12 #define PT\_YIELD(pt)

Yield from the current protothread.

This function will yield the protothread, thereby allowing other processing to take place in the system.

**Note:**

The [PT\\_YIELDING\(\)](#) flag must be placed first in the protothread's body if the [PT\\_YIELD\(\)](#) function should be used.

## Example

```
static
PT_THREAD(fade(struct pt *pt))
{
    PT_YIELDING();
    static int delay;

    PT_BEGIN(pt);

    for(delay = 3980; delay > 20; delay -= 20) {
        leds_red(LED_ON);
        clock_delay(4000 - delay);
        leds_red(LED_OFF);
        clock_delay(delay);
        PT_YIELD(pt);
    }

    PT_END(pt);
}
```

**Parameters:**

*pt* A pointer to the protothread control structure.

**4.1.2.13 #define PT\_YIELDING()**

Declare that a protothread can yield.

If a protothread should be able to yield with the [PT\\_YIELD\(\)](#) statement, this flag must be placed first in the protothread's function body.

## Example:

```
static
PT_THREAD(loop_thread(struct pt *pt))
{
    PT_YIELDING();
    static int i;

    PT_BEGIN(pt);

    for(i = 0; i < 200; ++i) {
        handle_item(i);
        PT_YIELD(pt);
    }

    PT_END(pt);
}
```

**4.2 Examples****4.2.1 A code-lock**

This example shows how to implement a simple code lock - the kind of device that is placed next to doors and that you have to push a four digit number into in order to unlock the door.

The code lock waits for key presses from a numeric keyboard and if the correct code is entered, the lock is unlocked. There is a maximum time of one second between each key press, and after the correct code has been entered, no more keys must be pressed for 0.5 seconds before the lock is opened.

```

/*
 * Copyright (c) 2004-2005, Swedish Institute of Computer Science.
 * All rights reserved.
 *
 * Redistribution and use in source and binary forms, with or without
 * modification, are permitted provided that the following conditions
 * are met:
 * 1. Redistributions of source code must retain the above copyright
 *    notice, this list of conditions and the following disclaimer.
 * 2. Redistributions in binary form must reproduce the above copyright
 *    notice, this list of conditions and the following disclaimer in the
 *    documentation and/or other materials provided with the distribution.
 * 3. Neither the name of the Institute nor the names of its contributors
 *    may be used to endorse or promote products derived from this software
 *    without specific prior written permission.
 *
 * THIS SOFTWARE IS PROVIDED BY THE INSTITUTE AND CONTRIBUTORS ``AS IS'' AND
 * ANY EXPRESS OR IMPLIED WARRANTIES, INCLUDING, BUT NOT LIMITED TO, THE
 * IMPLIED WARRANTIES OF MERCHANTABILITY AND FITNESS FOR A PARTICULAR PURPOSE
 * ARE DISCLAIMED. IN NO EVENT SHALL THE INSTITUTE OR CONTRIBUTORS BE LIABLE
 * FOR ANY DIRECT, INDIRECT, INCIDENTAL, SPECIAL, EXEMPLARY, OR CONSEQUENTIAL
 * DAMAGES (INCLUDING, BUT NOT LIMITED TO, PROCUREMENT OF SUBSTITUTE GOODS
 * OR SERVICES; LOSS OF USE, DATA, OR PROFITS; OR BUSINESS INTERRUPTION)
 * HOWEVER CAUSED AND ON ANY THEORY OF LIABILITY, WHETHER IN CONTRACT, STRICT
 * LIABILITY, OR TORT (INCLUDING NEGLIGENCE OR OTHERWISE) ARISING IN ANY WAY
 * OUT OF THE USE OF THIS SOFTWARE, EVEN IF ADVISED OF THE POSSIBILITY OF
 * SUCH DAMAGE.
 *
 * This file is part of the protothreads library.
 *
 * Author: Adam Dunkels <adam@sics.se>
 *
 * $Id: example-codelock.c,v 1.4 2005/04/01 09:20:39 adam Exp $
 */

/*
 *
 * This example shows how to implement a simple code lock. The code
 * lock waits for key presses from a numeric keyboard and if the
 * correct code is entered, the lock is unlocked. There is a maximum
 * time of one second between each key press, and after the correct
 * code has been entered, no more keys must be pressed for 0.5 seconds
 * before the lock is opened.
 *
 * This is an example that shows two things:
 * - how to implement a code lock key input mechanism, and
 * - how to implement a sequential timed routine.
 *
 * The program consists of two protothreads, one that implements the
 * code lock reader and one that implements simulated keyboard input.
 *
 */

#ifdef _WIN32
#include <windows.h>
#else
#include <unistd.h>
#include <sys/time.h>
#endif
#include <stdio.h>

#include "pt.h"

/*-----*/
/*
 * The following definitions are just for the simple timer library

```

```

* used in this example. The actual implementation of the functions
* can be found at the end of this file.
*/
struct timer { int start, interval; };
static int timer_expired(struct timer *t);
static void timer_set(struct timer *t, int usecs);
/*-----*/
/*
* This example uses two timers: one for the code lock protothread and
* one for the simulated key input protothread.
*/
static struct timer codelock_timer, input_timer;
/*-----*/
/*
* This is the code that has to be entered.
*/
static const char code[4] = {'1', '4', '2', '3'};
/*-----*/
/*
* This example has two protothread and therefor has two protothread
* control structures of type struct pt. These are initialized with
* PT_INIT() in the main() function below.
*/
static struct pt codelock_pt, input_pt;
/*-----*/
/*
* The following code implements a simple key input. Input is made
* with the press_key() function, and the function key_pressed()
* checks if a key has been pressed. The variable "key" holds the
* latest key that was pressed. The variable "key_pressed_flag" is set
* when a key is pressed and cleared when a key press is checked.
*/
static char key, key_pressed_flag;

static void
press_key(char k)
{
    printf("--- Key '%c' pressed\n", k);
    key = k;
    key_pressed_flag = 1;
}

static int
key_pressed(void)
{
    if(key_pressed_flag != 0) {
        key_pressed_flag = 0;
        return 1;
    }
    return 0;
}
/*-----*/
/*
* Declaration of the protothread function implementing the code lock
* logic. The protothread function is declared using the PT_THREAD()
* macro. The function is declared with the "static" keyword since it
* is local to this file. The name of the function is codelock_thread
* and it takes one argument, pt, of the type struct pt.
*
*/
static
PT_THREAD(codelock_thread(struct pt *pt))
{
    /* This is a local variable that holds the number of keys that have
    * been pressed. Note that it is declared with the "static" keyword
    * to make sure that the variable is *not* allocated on the stack.
    */

```

```

static int keys;

/*
 * Declare the beginning of the protothread.
 */
PT_BEGIN(pt);

/*
 * We'll let the protothread loop until the protothread is
 * explicitly exited with PT_EXIT().
 */
while(1) {

    /*
     * We'll be reading key presses until we get the right amount of
     * correct keys.
     */
    for(keys = 0; keys < sizeof(code); ++keys) {

        /*
         * If we haven't gotten any keypresses, we'll simply wait for one.
         */
        if(keys == 0) {

            /*
             * The PT_WAIT_UNTIL() function will block until the condition
             * key_pressed() is true.
             */
            PT_WAIT_UNTIL(pt, key_pressed());
        } else {

            /*
             * If the "key" variable was larger than zero, we have already
             * gotten at least one correct key press. If so, we'll not
             * only wait for the next key, but we'll also set a timer that
             * expires in one second. This gives the person pressing the
             * keys one second to press the next key in the code.
             */
            timer_set(&codelock_timer, 1000);

            /*
             * The following statement shows how complex blocking
             * conditions can be easily expressed with protothreads and
             * the PT_WAIT_UNTIL() function.
             */
            PT_WAIT_UNTIL(pt, key_pressed() || timer_expired(&codelock_timer));

            /*
             * If the timer expired, we should break out of the for() loop
             * and start reading keys from the beginning of the while(1)
             * loop instead.
             */
            if(timer_expired(&codelock_timer)) {
                printf("Code lock timer expired.\n");

                /*
                 * Break out from the for() loop and start from the
                 * beginning of the while(1) loop.
                 */
                break;
            }
        }
    }

    /*
     * Check if the pressed key was correct.
     */
    if(key != code[keys]) {

```

```

        printf("Incorrect key '%c' found\n", key);
        /*
         * Break out of the for() loop since the key was incorrect.
         */
        break;
    } else {
        printf("Correct key '%c' found\n", key);
    }
}

/*
 * Check if we have gotten all keys.
 */
if(keys == sizeof(code)) {
    printf("Correct code entered, waiting for 500 ms before unlocking.\n");

    /*
     * Ok, we got the correct code. But to make sure that the code
     * was not just a fluke of luck by an intruder, but the correct
     * code entered by a person that knows the correct code, we'll
     * wait for half a second before opening the lock. If another
     * key is pressed during this time, we'll assume that it was a
     * fluke of luck that the correct code was entered the first
     * time.
     */
    timer_set(&codelock_timer, 500);
    PT_WAIT_UNTIL(pt, key_pressed() || timer_expired(&codelock_timer));

    /*
     * If we continued from the PT_WAIT_UNTIL() statement without
     * the timer expired, we don't open the lock.
     */
    if(!timer_expired(&codelock_timer)) {
        printf("Key pressed during final wait, code lock locked again.\n");
    } else {

        /*
         * If the timer expired, we'll open the lock and exit from the
         * protothread.
         */
        printf("Code lock unlocked.\n");
        PT_EXIT(pt);
    }
}

/*
 * Finally, we'll mark the end of the protothread.
 */
PT_END(pt);
}

/*-----*/
/*
 * This is the second protothread in this example. It implements a
 * simulated user pressing the keys. This illustrates how a linear
 * sequence of timed instructions can be implemented with
 * protothreads.
 */
static
PT_THREAD(input_thread(struct pt *pt))
{
    PT_BEGIN(pt);

    printf("Waiting 1 second before entering first key.\n");

    timer_set(&input_timer, 1000);
    PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

```

```
press_key('1');

timer_set(&input_timer, 100);
PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

press_key('2');

timer_set(&input_timer, 100);
PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

press_key('3');

timer_set(&input_timer, 2000);
PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

press_key('1');

timer_set(&input_timer, 200);
PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

press_key('4');

timer_set(&input_timer, 200);
PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

press_key('2');

timer_set(&input_timer, 2000);
PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

press_key('3');

timer_set(&input_timer, 200);
PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

press_key('1');

timer_set(&input_timer, 200);
PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

press_key('4');

timer_set(&input_timer, 200);
PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

press_key('2');

timer_set(&input_timer, 100);
PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

press_key('3');

timer_set(&input_timer, 100);
PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

press_key('4');

timer_set(&input_timer, 1500);
PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

press_key('1');

timer_set(&input_timer, 300);
PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

press_key('4');
```



```

    timer_set(&input_timer, 400);
    PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

    press_key('2');

    timer_set(&input_timer, 500);
    PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

    press_key('3');

    timer_set(&input_timer, 2000);
    PT_WAIT_UNTIL(pt, timer_expired(&input_timer));

    PT_END(pt);
}
/*-----*/
/*
 * This is the main function. It initializes the two protothread
 * control structures and schedules the two protothreads. The main
 * function returns when the protothread the runs the code lock exits.
 */
int
main(void)
{
    /*
     * Initialize the two protothread control structures.
     */
    PT_INIT(&input_pt);
    PT_INIT(&codelock_pt);

    /*
     * Schedule the two protothreads until the codelock_thread() exits.
     */
    while(PT_SCHEDULE(codelock_thread(&codelock_pt))) {
        PT_SCHEDULE(input_thread(&input_pt));

        /*
         * When running this example on a multitasking system, we must
         * give other processes a chance to run too and therefore we call
         * usleep() resp. Sleep() here. On a dedicated embedded system,
         * we usually do not need to do this.
         */
#ifdef _WIN32
        Sleep(0);
#else
        usleep(10);
#endif
    }

    return 0;
}
/*-----*/
/*
 * Finally, the implementation of the simple timer library follows.
 */
#ifdef _WIN32

static int clock_time(void)
{ return (int)GetTickCount(); }

#else /* _WIN32 */

static int clock_time(void)
{
    struct timeval tv;
    struct timezone tz;

```

```

    gettimeofday(&tv, &tz);
    return tv.tv_sec * 1000 + tv.tv_usec / 1000;
}

#endif /* _WIN32 */

static int timer_expired(struct timer *t)
{ return (int)(clock_time() - t->start) >= (int)t->interval; }

static void timer_set(struct timer *t, int interval)
{ t->interval = interval; t->start = clock_time(); }
/*-----*/

```

### 4.2.2 The bounded buffer with protothread semaphores

The following example shows how to implement the bounded buffer problem using the protothreads semaphore library. The example uses three protothreads: one `producer()` protothread that produces items, one `consumer()` protothread that consumes items, and one `driver_thread()` that schedules the producer and consumer protothreads.

Note that there is no need for a mutex to guard the `add_to_buffer()` and `get_from_buffer()` functions because of the implicit locking semantics of protothreads - a protothread will never be preempted and will never block except in an explicit `PT_WAIT` statement.

```

/*
 * Copyright (c) 2004-2005, Swedish Institute of Computer Science.
 * All rights reserved.
 *
 * Redistribution and use in source and binary forms, with or without
 * modification, are permitted provided that the following conditions
 * are met:
 * 1. Redistributions of source code must retain the above copyright
 *    notice, this list of conditions and the following disclaimer.
 * 2. Redistributions in binary form must reproduce the above copyright
 *    notice, this list of conditions and the following disclaimer in the
 *    documentation and/or other materials provided with the distribution.
 * 3. Neither the name of the Institute nor the names of its contributors
 *    may be used to endorse or promote products derived from this software
 *    without specific prior written permission.
 *
 * THIS SOFTWARE IS PROVIDED BY THE INSTITUTE AND CONTRIBUTORS ``AS IS'' AND
 * ANY EXPRESS OR IMPLIED WARRANTIES, INCLUDING, BUT NOT LIMITED TO, THE
 * IMPLIED WARRANTIES OF MERCHANTABILITY AND FITNESS FOR A PARTICULAR PURPOSE
 * ARE DISCLAIMED. IN NO EVENT SHALL THE INSTITUTE OR CONTRIBUTORS BE LIABLE
 * FOR ANY DIRECT, INDIRECT, INCIDENTAL, SPECIAL, EXEMPLARY, OR CONSEQUENTIAL
 * DAMAGES (INCLUDING, BUT NOT LIMITED TO, PROCUREMENT OF SUBSTITUTE GOODS
 * OR SERVICES; LOSS OF USE, DATA, OR PROFITS; OR BUSINESS INTERRUPTION)
 * HOWEVER CAUSED AND ON ANY THEORY OF LIABILITY, WHETHER IN CONTRACT, STRICT
 * LIABILITY, OR TORT (INCLUDING NEGLIGENCE OR OTHERWISE) ARISING IN ANY WAY
 * OUT OF THE USE OF THIS SOFTWARE, EVEN IF ADVISED OF THE POSSIBILITY OF
 * SUCH DAMAGE.
 *
 * This file is part of the protothreads library.
 *
 * Author: Adam Dunkels <adam@sics.se>
 *
 * $Id: example-buffer.c,v 1.3 2005/04/01 09:20:39 adam Exp $
 */

#ifdef _WIN32
#include <windows.h>
#else
#include <unistd.h>
#endif

```

```
#include <stdio.h>

#include "pt-sem.h"

#define NUM_ITEMS 32
#define BUFSIZE 8

static int buffer[BUFSIZE];
static int bufptr;

static void
add_to_buffer(int item)
{
    printf("Item %d added to buffer at place %d\n", item, bufptr);
    buffer[bufptr] = item;
    bufptr = (bufptr + 1) % BUFSIZE;
}

static int
get_from_buffer(void)
{
    int item;
    item = buffer[bufptr];
    printf("Item %d retrieved from buffer at place %d\n",
           item, bufptr);
    bufptr = (bufptr + 1) % BUFSIZE;
    return item;
}

static int
produce_item(void)
{
    static int item = 0;
    printf("Item %d produced\n", item);
    return item++;
}

static void
consume_item(int item)
{
    printf("Item %d consumed\n", item);
}

static struct pt_sem full, empty;

static
PT_THREAD(producer(struct pt *pt))
{
    static int produced;

    PT_BEGIN(pt);

    for(produced = 0; produced < NUM_ITEMS; ++produced) {

        PT_SEM_WAIT(pt, &full);

        add_to_buffer(produce_item());

        PT_SEM_SIGNAL(pt, &empty);
    }

    PT_END(pt);
}

static
PT_THREAD(consumer(struct pt *pt))
{
    static int consumed;
```

```

PT_BEGIN(pt);

for(consumed = 0; consumed < NUM_ITEMS; ++consumed) {

    PT_SEM_WAIT(pt, &empty);

    consume_item(get_from_buffer());

    PT_SEM_SIGNAL(pt, &full);
}

PT_END(pt);
}

static
PT_THREAD(driver_thread(struct pt *pt))
{
    static struct pt pt_producer, pt_consumer;

    PT_BEGIN(pt);

    PT_SEM_INIT(&empty, 0);
    PT_SEM_INIT(&full, BUFSIZE);

    PT_INIT(&pt_producer);
    PT_INIT(&pt_consumer);

    PT_WAIT_THREAD(pt, producer(&pt_producer) &
                    consumer(&pt_consumer));

    PT_END(pt);
}

int
main(void)
{
    struct pt driver_pt;

    PT_INIT(&driver_pt);

    while(PT_SCHEDULE(driver_thread(&driver_pt))) {

        /*
         * When running this example on a multitasking system, we must
         * give other processes a chance to run too and therefore we call
         * usleep() resp. Sleep() here. On a dedicated embedded system,
         * we usually do not need to do this.
         */
#ifdef _WIN32
        Sleep(0);
#else
        usleep(10);
#endif
    }
    return 0;
}

```

## 4.3 Protothread semaphores

### 4.3.1 Detailed Description

This module implements counting semaphores on top of protothreads. Semaphores are a synchronization primitive that provide two operations: "wait" and "signal". The "wait" operation checks the semaphore

counter and blocks the thread if the counter is zero. The "signal" operation increases the semaphore counter but does not block. If another thread has blocked waiting for the semaphore that is signalled, the blocked thread will become runnable again.

Semaphores can be used to implement other, more structured, synchronization primitives such as monitors and message queues/bounded buffers (see below).

The following example shows how the producer-consumer problem, also known as the bounded buffer problem, can be solved using protothreads and semaphores. Notes on the program follow after the example.

```
#include "pt-sem.h"

#define NUM_ITEMS 32
#define BUFSIZE 8

static struct pt_sem mutex, full, empty;

PT_THREAD(producer(struct pt *pt))
{
    static int produced;

    PT_BEGIN(pt);

    for(produced = 0; produced < NUM_ITEMS; ++produced) {

        PT_SEM_WAIT(pt, &full);

        PT_SEM_WAIT(pt, &mutex);
        add_to_buffer(produce_item());
        PT_SEM_SIGNAL(pt, &mutex);

        PT_SEM_SIGNAL(pt, &empty);
    }

    PT_END(pt);
}

PT_THREAD(consumer(struct pt *pt))
{
    static int consumed;

    PT_BEGIN(pt);

    for(consumed = 0; consumed < NUM_ITEMS; ++consumed) {

        PT_SEM_WAIT(pt, &empty);

        PT_SEM_WAIT(pt, &mutex);
        consume_item(get_from_buffer());
        PT_SEM_SIGNAL(pt, &mutex);

        PT_SEM_SIGNAL(pt, &full);
    }

    PT_END(pt);
}

PT_THREAD(driver_thread(struct pt *pt))
{
    static struct pt pt_producer, pt_consumer;

    PT_BEGIN(pt);

    PT_SEM_INIT(&empty, 0);
    PT_SEM_INIT(&full, BUFSIZE);
    PT_SEM_INIT(&mutex, 1);
```

```

PT_INIT(&pt_producer);
PT_INIT(&pt_consumer);

PT_WAIT_THREAD(pt, producer(&pt_producer) &
               consumer(&pt_consumer));

PT_END(pt);
}

```

The program uses three protothreads: one protothread that implements the consumer, one thread that implements the producer, and one protothread that drives the two other protothreads. The program uses three semaphores: "full", "empty" and "mutex". The "mutex" semaphore is used to provide mutual exclusion for the buffer, the "empty" semaphore is used to block the consumer if the buffer is empty, and the "full" semaphore is used to block the producer if the buffer is full.

The "driver\_thread" holds two protothread state variables, "pt\_producer" and "pt\_consumer". It is important to note that both these variables are declared as *static*. If the static keyword is not used, both variables are stored on the stack. Since protothreads do not store the stack, these variables may be overwritten during a protothread wait operation. Similarly, both the "consumer" and "producer" protothreads declare their local variables as static, to avoid them being stored on the stack.

## Files

- file [pt-sem.h](#)  
*Couting semaphores implemented on protothreads.*

## Defines

- #define [PT\\_SEM\\_INIT](#)(s, c)  
*Initialize a semaphore.*
- #define [PT\\_SEM\\_WAIT](#)(pt, s)  
*Wait for a semaphore.*
- #define [PT\\_SEM\\_SIGNAL](#)(pt, s)  
*Signal a semaphore.*

### 4.3.2 Define Documentation

#### 4.3.2.1 #define PT\_SEM\_INIT(s, c)

Initialize a semaphore.

This macro initializes a semaphore with a value for the counter. Internally, the semaphores use an "unsigned int" to represent the counter, and therefore the "count" argument should be within range of an unsigned int.

#### Parameters:

- s* (struct pt\_sem \*) A pointer to the pt\_sem struct representing the semaphore
- c* (unsigned int) The initial count of the semaphore.

#### 4.3.2.2 #define PT\_SEM\_SIGNAL(pt, s)

Signal a semaphore.

This macro carries out the "signal" operation on the semaphore. The signal operation increments the counter inside the semaphore, which eventually will cause waiting protothreads to continue executing.

**Parameters:**

- pt* (struct pt \*) A pointer to the protothread (struct pt) in which the operation is executed.
- s* (struct pt\_sem \*) A pointer to the pt\_sem struct representing the semaphore

#### 4.3.2.3 #define PT\_SEM\_WAIT(pt, s)

Wait for a semaphore.

This macro carries out the "wait" operation on the semaphore. The wait operation causes the protothread to block while the counter is zero. When the counter reaches a value larger than zero, the protothread will continue.

**Parameters:**

- pt* (struct pt \*) A pointer to the protothread (struct pt) in which the operation is executed.
- s* (struct pt\_sem \*) A pointer to the pt\_sem struct representing the semaphore

## 4.4 Local continuations

### 4.4.1 Detailed Description

Local continuations form the basis for implementing protothreads. A local continuation can be *set* in a specific function to capture the state of the function. After a local continuation has been set can be *resumed* in order to restore the state of the function at the point where the local continuation was set.

#### Files

- file [lc.h](#)  
*Local continuations.*
- file [lc-switch.h](#)  
*Implementation of local continuations based on switch() statment.*
- file [lc-addrlabels.h](#)  
*Implementation of local continuations based on the "Labels as values" feature of gcc.*

#### Defines

- #define [LC\\_INIT](#)(lc)  
*Initialize a local continuation.*
- #define [LC\\_SET](#)(lc)  
*Set a local continuation.*

- `#define LC_RESUME(lc)`  
*Resume a local continuation.*
- `#define LC_END(lc)`  
*Mark the end of local continuation usage.*

## Typedefs

- `typedef unsigned short lc_t`  
*The local continuation type.*

### 4.4.2 Define Documentation

#### 4.4.2.1 `#define LC_END(lc)`

Mark the end of local continuation usage.

The end operation signifies that local continuations should not be used any more in the function. This operation is not needed for most implementations of local continuation, but is required by a few implementations.

#### 4.4.2.2 `#define LC_INIT(lc)`

Initialize a local continuation.

This operation initializes the local continuation, thereby unsetting any previously set continuation state.

#### 4.4.2.3 `#define LC_RESUME(lc)`

Resume a local continuation.

The resume operation resumes a previously set local continuation, thus restoring the state in which the function was when the local continuation was set. If the local continuation has not been previously set, the resume operation does nothing.

#### 4.4.2.4 `#define LC_SET(lc)`

Set a local continuation.

The set operation saves the state of the function at the point where the operation is executed. As far as the set operation is concerned, the state of the function does **not** include the call-stack or local (automatic) variables, but only the program counter and such CPU registers that needs to be saved.

## 5 The Protothreads Library 1.1 File Documentation

### 5.1 `lc-addrlabels.h` File Reference

#### 5.1.1 Detailed Description

Implementation of local continuations based on the "Labels as values" feature of gcc.



**Author:**

Adam Dunkels <[adam@sics.se](mailto:adam@sics.se)>

This implementation of local continuations is based on a special feature of the GCC C compiler called "labels as values". This feature allows assigning pointers with the address of the code corresponding to a particular C label.

For more information, see the GCC documentation: <http://gcc.gnu.org/onlinedocs/gcc/Labels-as-Values>

## 5.2 lc-switch.h File Reference

### 5.2.1 Detailed Description

Implementation of local continuations based on switch() statment.

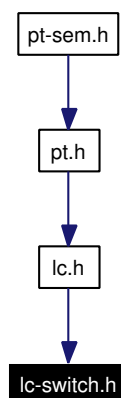
**Author:**

Adam Dunkels <[adam@sics.se](mailto:adam@sics.se)>

This implementation of local continuations uses the C switch() statement to resume execution of a function somewhere inside the function's body. The implementation is based on the fact that switch() statements are able to jump directly into the bodies of control structures such as if() or while() statmenets.

This implementation borrows heavily from Simon Tatham's coroutines implementation in C: <http://www.chiark.greenend.org.uk/~sgtatham/coroutines.html>

This graph shows which files directly or indirectly include this file:

**Typedefs**

- typedef unsigned short [lc\\_t](#)  
*The local continuation type.*

## 5.3 lc.h File Reference

### 5.3.1 Detailed Description

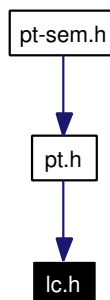
Local continuations.

**Author:**

Adam Dunkels <adam@sics.se>

```
#include "lc-switch.h"
```

This graph shows which files directly or indirectly include this file:

**Defines**

- #define `LC_INIT(lc)`  
*Initialize a local continuation.*
- #define `LC_SET(lc)`  
*Set a local continuation.*
- #define `LC_RESUME(lc)`  
*Resume a local continuation.*
- #define `LC_END(lc)`  
*Mark the end of local continuation usage.*

## 5.4 pt-sem.h File Reference

### 5.4.1 Detailed Description

Couting semaphores implemented on protothreads.

**Author:**

Adam Dunkels <adam@sics.se>

```
#include "pt.h"
```

**Defines**

- #define `PT_SEM_INIT(s, c)`  
*Initialize a semaphore.*

- #define `PT_SEM_WAIT`(pt, s)  
*Wait for a semaphore.*
- #define `PT_SEM_SIGNAL`(pt, s)  
*Signal a semaphore.*

## 5.5 pt.h File Reference

### 5.5.1 Detailed Description

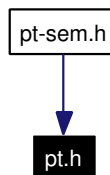
Protothreads implementation.

**Author:**

Adam Dunkels <[adam@sics.se](mailto:adam@sics.se)>

```
#include "lc.h"
```

This graph shows which files directly or indirectly include this file:



### Defines

- #define `PT_THREAD`(name\_args)  
*Declaration of a protothread.*
- #define `PT_INIT`(pt)  
*Initialize a protothread.*
- #define `PT_BEGIN`(pt)  
*Declare the start of a protothread inside the C function implementing the protothread.*
- #define `PT_WAIT_UNTIL`(pt, condition)  
*Block and wait until condition is true.*
- #define `PT_WAIT_WHILE`(pt, cond)  
*Block and wait while condition is true.*
- #define `PT_WAIT_THREAD`(pt, thread)  
*Block and wait until a child protothread completes.*
- #define `PT_SPAWN`(pt, child, thread)  
*Spawn a child protothread and wait until it exits.*

- #define `PT_RESTART`(pt)  
*Restart the protothread.*
- #define `PT_EXIT`(pt)  
*Exit the protothread.*
- #define `PT_END`(pt)  
*Declare the end of a protothread.*
- #define `PT_SCHEDULE`(f)  
*Schedule a protothread.*
- #define `PT_YIELDING`()  
*Declarare that a protothread can yield.*
- #define `PT_YIELD`(pt)  
*Yield from the current protothread.*

## Index

Examples, [10](#)

lc

- LC\_END, [23](#)
- LC\_INIT, [23](#)
- LC\_RESUME, [23](#)
- LC\_SET, [23](#)

lc-addrlabels.h, [23](#)

lc-switch.h, [24](#)

lc.h, [24](#)

LC\_END  
lc, [23](#)

LC\_INIT  
lc, [23](#)

LC\_RESUME  
lc, [23](#)

LC\_SET  
lc, [23](#)

Local continuations, [22](#)

Protothread semaphores, [19](#)

Protothreads, [4](#)

pt

- PT\_BEGIN, [5](#)
- PT\_END, [6](#)
- PT\_EXIT, [6](#)
- PT\_INIT, [6](#)
- PT\_RESTART, [6](#)
- PT\_SCHEDULE, [7](#)
- PT\_SPAWN, [7](#)
- PT\_THREAD, [8](#)
- PT\_WAIT\_THREAD, [8](#)
- PT\_WAIT\_UNTIL, [9](#)
- PT\_WAIT\_WHILE, [9](#)
- PT\_YIELD, [9](#)
- PT\_YIELDING, [10](#)

pt-sem.h, [25](#)

pt.h, [26](#)

PT\_BEGIN  
pt, [5](#)

PT\_END  
pt, [6](#)

PT\_EXIT  
pt, [6](#)

PT\_INIT  
pt, [6](#)

PT\_RESTART  
pt, [6](#)

PT\_SCHEDULE  
pt, [7](#)

PT\_SEM\_INIT

ptsem, [21](#)

PT\_SEM\_SIGNAL  
ptsem, [21](#)

PT\_SEM\_WAIT  
ptsem, [22](#)

PT\_SPAWN  
pt, [7](#)

PT\_THREAD  
pt, [8](#)

PT\_WAIT\_THREAD  
pt, [8](#)

PT\_WAIT\_UNTIL  
pt, [9](#)

PT\_WAIT\_WHILE  
pt, [9](#)

PT\_YIELD  
pt, [9](#)

PT\_YIELDING  
pt, [10](#)

ptsem

PT\_SEM\_INIT, [21](#)

PT\_SEM\_SIGNAL, [21](#)

PT\_SEM\_WAIT, [22](#)